Structured Polynomial Codes

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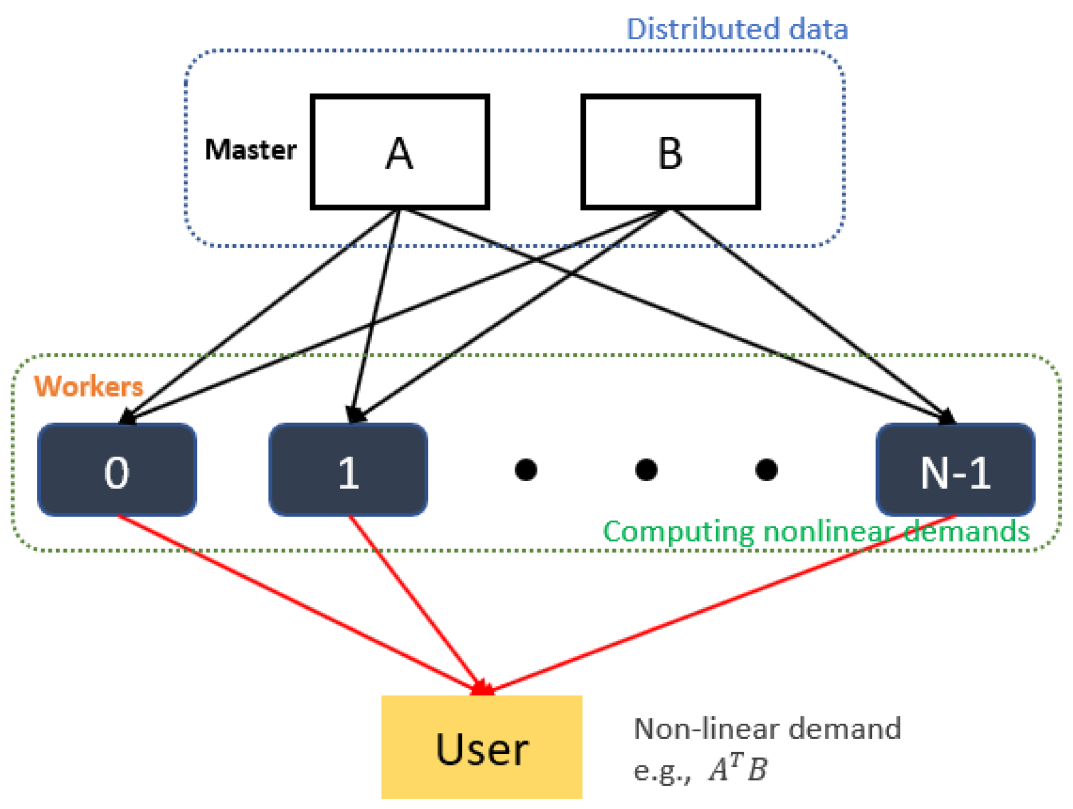
Motivation



- A fundamental challenge: Balancing computation and communication complexity.

Related work

- **Frameworks for distributed computing:** MapReduce, Hadoop, Spark, TeraSort [1]
- **Channel coding approaches:** Polynomial codes, Lagrange coded computing [2, 3]
- **Source coding approaches:** Structured codes for modulo two sum computation in [4], and distributed matrix multiplication in [5]

Contributions

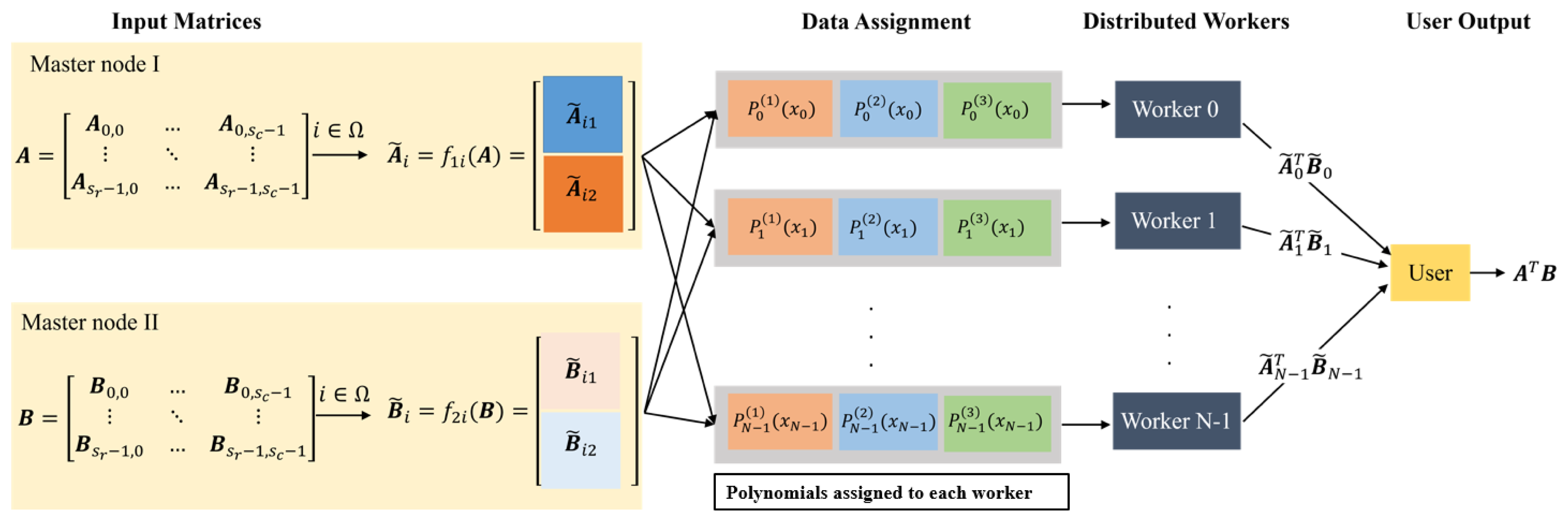
Novelty:

- Combining the benefits of structured coding and polynomial codes
- Elevating the Körner-Marton approach to the distributed matrix multiplication setting
- Incorporating a secure matrix multiplication design

Savings:

- low complexity distributed encoding
- communication costs (reduced by %50)
- storage size (reduced by %50)

A structured distributed matrix multiplication model



- ☐ *Each worker, using the assigned polynomials, calculates the product of sub-matrices $\tilde{\mathbf{A}}_i^\intercal \tilde{\mathbf{B}}_i$.*
- ☐ *Using $\{\tilde{\mathbf{A}}_i^\intercal \tilde{\mathbf{B}}_i\}_i$ from a subset of workers, the user decodes $\mathbf{AB}$.*
- ☐ *The user cannot decode $\mathbf{A}$ or $\mathbf{B}$, where the security of multiplication is ensured by structured coding.*

Source coding for matrix multiplication [5]

Two *distributed sources*, $\mathbf{A} \in \mathbb{F}_q^{m\times 1}$ and $\mathbf{B} \in \mathbb{F}_q^{m\times 1}$:

- Splitting of each source: $\mathbf{A} = \begin{bmatrix} \mathbf{A}_1 \\ \mathbf{A}_2 \end{bmatrix}^\intercal \in \mathbb{F}_q^{m\times 1}$, $\mathbf{B} = \begin{bmatrix} \mathbf{B}_1 \\ \mathbf{B}_2 \end{bmatrix} \in \mathbb{F}_q^{m\times 1}$,
- Nonlinear mapping from each source:

$$\mathbf{X}_1 = g_1(\mathbf{A}) = \begin{bmatrix} \mathbf{A}_2 \\ \mathbf{A}_1 \\ \mathbf{A}_2^\intercal \mathbf{A}_1 \end{bmatrix} \in \mathbb{F}_2^{(m+1)\times 1} , \qquad \mathbf{X}_2 = g_2(\mathbf{B}) = \begin{bmatrix} \mathbf{B}_1 \\ \mathbf{B}_2 \\ \mathbf{B}_1^\intercal \mathbf{B}_2 \end{bmatrix} \in \mathbb{F}_2^{(m+1)\times 1} .$$

- Linear encoding: Sources use a common encoder, and compute $\mathbf{CX}_j^n \in \mathbb{F}_2^{(m+1)\times k}$ and send $\mathbf{CX}_j^n$ [4].
- Decoding: Exploiting [4], the sum rate needed for the user to recover the vector sequence

$$\mathbf{Z}^n = \mathbf{X}_1^n \oplus_2 \mathbf{X}_2^n \in \mathbb{F}_2^{(m+1)\times n}$$

with a vanishing error probability, is determined as:

$$R_{\mathrm{KM}}^{\Sigma} = 2H(\mathbf{X}_1 \oplus_2 \mathbf{X}_2) = 2H(U, V, W) ,$$

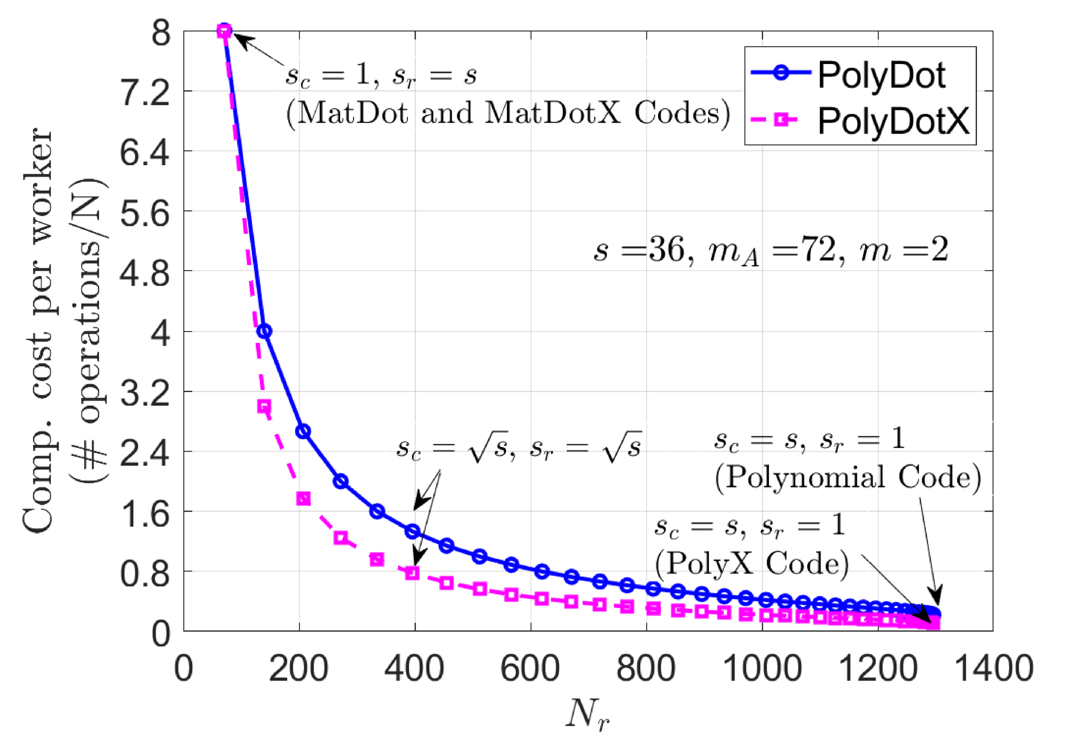
where the following vectors can be computed in a fully distributed manner:

$$\mathbf{U} = A_2 \oplus_q B_1 \in \mathbb{F}_q^{m/2\times 1} , \quad \mathbf{V} = A_1 \oplus_q B_2 \in \mathbb{F}_q^{m/2\times 1} , \quad W = A_2^T A_1 \oplus_q B_1^T B_2 \in \mathbb{F}_q .$$

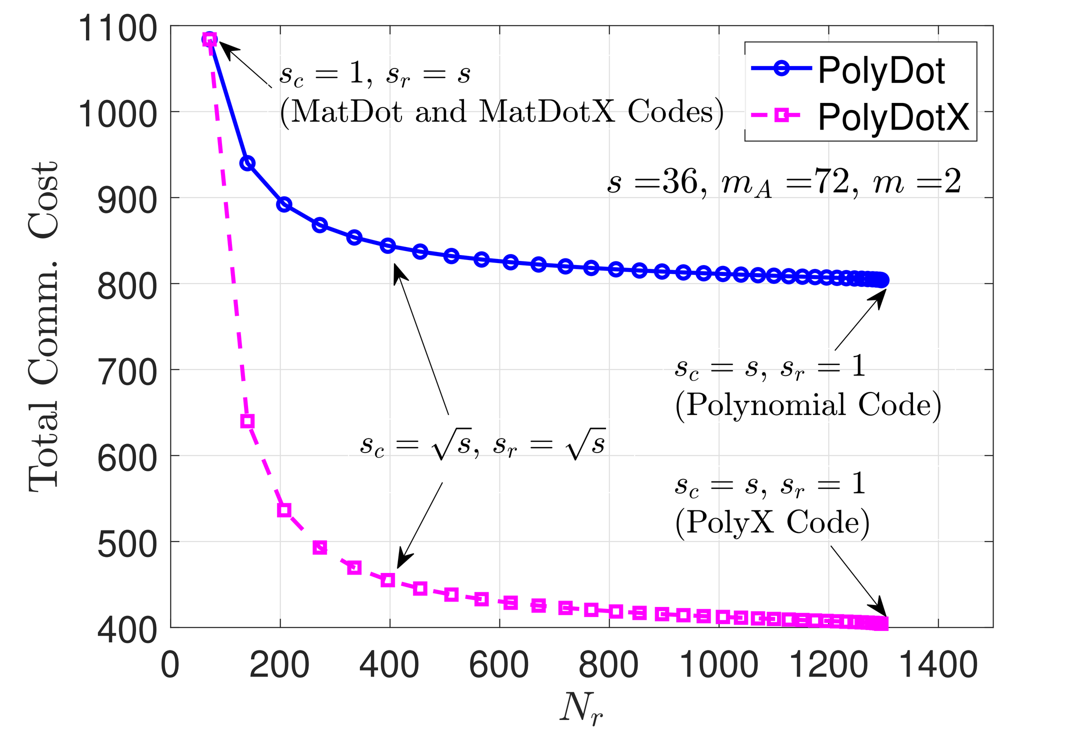
The user can recover the desired inner product using $\mathbf{U}$, $\mathbf{V}$, and W.

Performance results

For $s_c \gg m$, the upper bound of computation cost per worker approaches $1 + \frac{1}{2s}$.



The total communication cost is reduced by %50 compared to the PolyDot model.



Future directions

Structured codes for

- n-matrix products
- privacy/security aspects
- tensor product computations

References

[1] Alkatheri et al. A comparative study of big data frameworks. *Int. Jour. Comp. Sci. IJCSIS*, 2019.

[2] López et al. Secure MatDot codes: a secure, distributed matrix multiplication scheme. In *ITW 2022*, Mumbai, India, 2022.

[3] Yu et al. Lagrange coded computing: Optimal design for resiliency, security, and privacy. In *Proc. Int. on AI and Stat.*, 2019.

[4] Körner and Marton. How to encode the modulo-two sum of binary sources. *IEEE Trans. Inf. Theory*, 1979.

[5] Malak. Distributed structured matrix multiplication. In *ISIT*, Athens, Greece, Jul. 2024.